

Data Structures and Algorithms

Prepared by Prof. Shereen Taie

IT251
Information Technology Department

Lecture 4



Data Structures and Algorithms

Lecture 4

IT251

Information Technology Department

Prepared by
Prof. Shereen Taie

www.midocean.ae

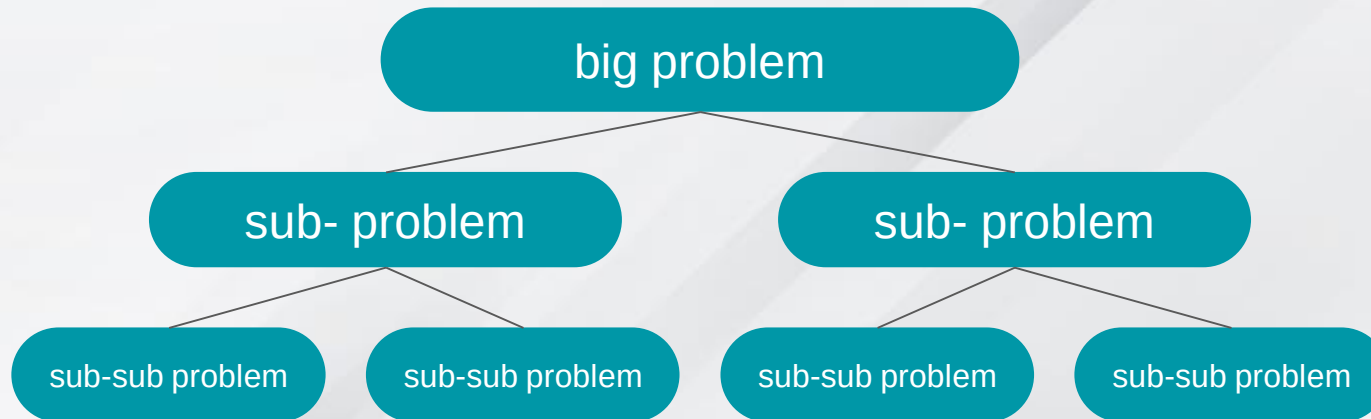
Contents:

- Divide And Conquer
 - Divide-and-Conquer Examples
- Mergesort
 - Mergesort: Pseudocode
 - Mergesort: Recursive Calls
 - Mergesort: Merge Steps
 - Mergesort: Is It Fast?
 - MERGESORT: $O(n \log n)$ PROOF
 - Mergesort Recursion Tree
- Recurrence Relations
 - Recurrence Relations Examples

DIVIDE AND CONQUER (Merge Sort - Recurrence Relations)

DIVIDE AND CONQUER

- **An algorithm design paradigm:**
 1. break up a problem into smaller subproblems
 2. solve those subproblems *recursively*
 3. combine the results of those subproblems to get the overall answer



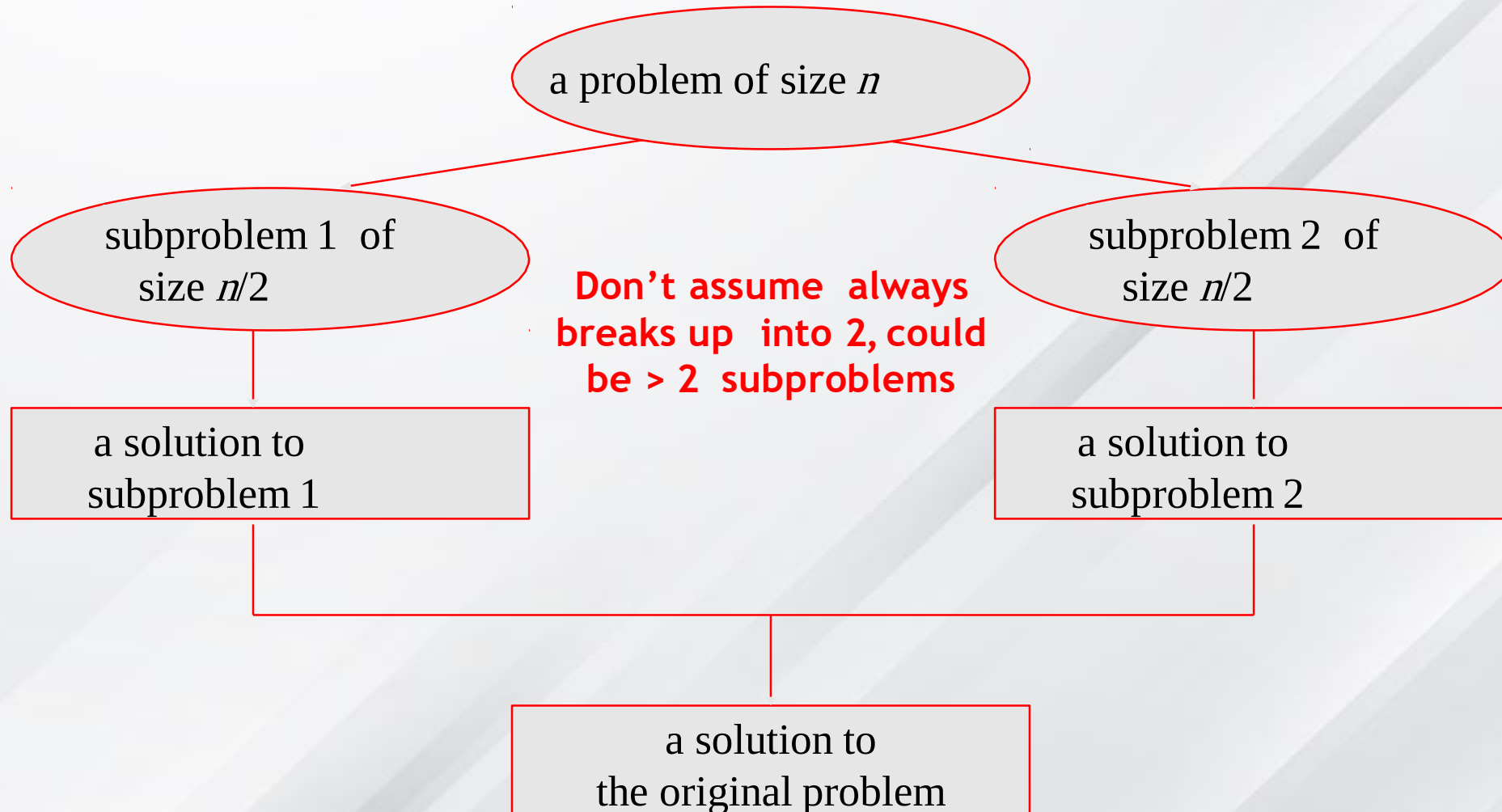
Divide-and-Conquer Examples

- Sorting: mergesort and quicksort
- Binary tree traversals
- Binary search
- Multiplication of large integers
- Matrix multiplication: Strassen's algorithm

Divide and conquer

- Three steps are involved:
 - ✓ Divide the problem into a number of sub-problems
 - ✓ Conquer the sub-problems (solve them)
 - ✓ Combine the sub-problem solutions to get the solution to the original problem
- Note: often the “conquer” step is done recursively.

Divide-and-conquer technique



Divide and Conquer

Based on dividing problem into sub-problems

Approach

- Divide problem into smaller sub-problems
 1. Sub-problems must be of same type
 2. Sub-problems do not need to overlap
- Conquer by solving sub-problems recursively.
- If the sub-problems are small enough, solve them in brute force fashion
- Combine the solutions of sub-problems into a solution of the original problem.
- Usually contains two or more recursive calls.

Divide-and-conquer Algorithms

- For Divide-and-Conquer algorithms the running time is mainly affected by 3 criteria:

The number of sub-problems

The ratio of initial problem size to sub- problem size.

The number of steps required to:

divide the initial problem

combine sub-solutions.

Analyzing Divide-and Conquer Algorithms

The master theorem applies to recursive algorithms of the form:

$$T(n) = a T(n/b) + f(n),$$

Where,

n is the size of the input,

a is the number of recursive calls,

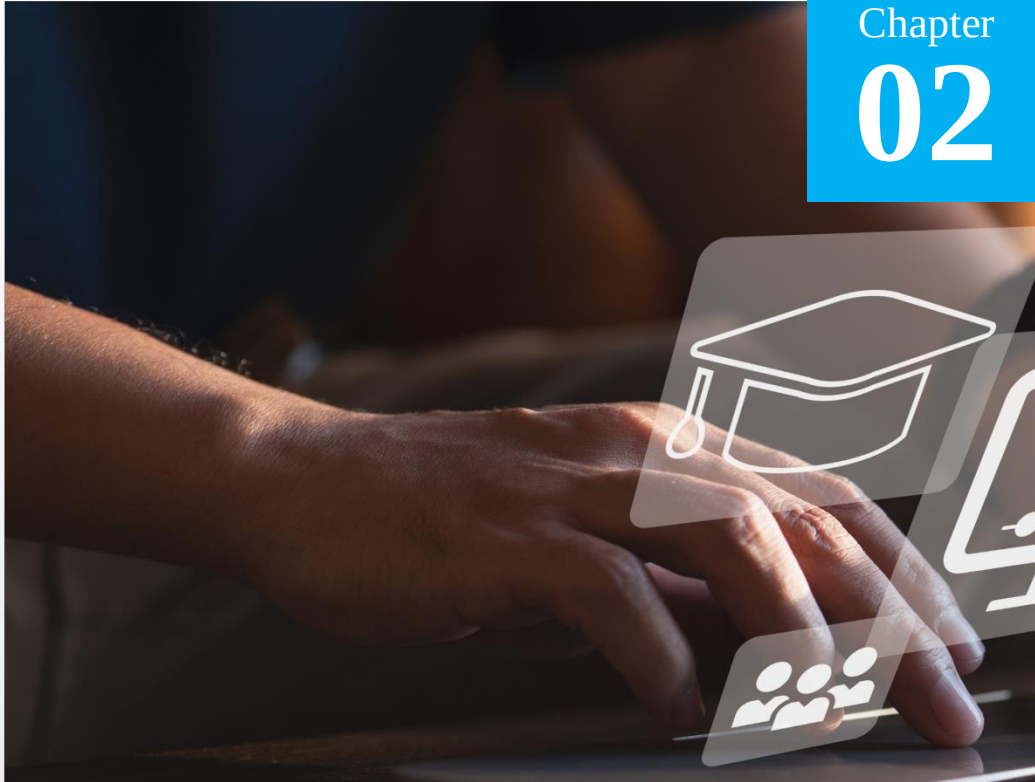
b is the factor by which the input is divided, and

f(n) is the non-recursive work done at each level (time to divide the problem of size n, time to combine the sub-problem solutions).

Chapter
02

Divide-and Conquer Example

Merge Sort



Merge sort:

- Merge sort is a sorting technique based on divide and conquer technique.
- It has a worst-case time complexity of $O(n \log n)$.
- Merge sort first divides the array into equal halves and then combines them in a sorted manner.

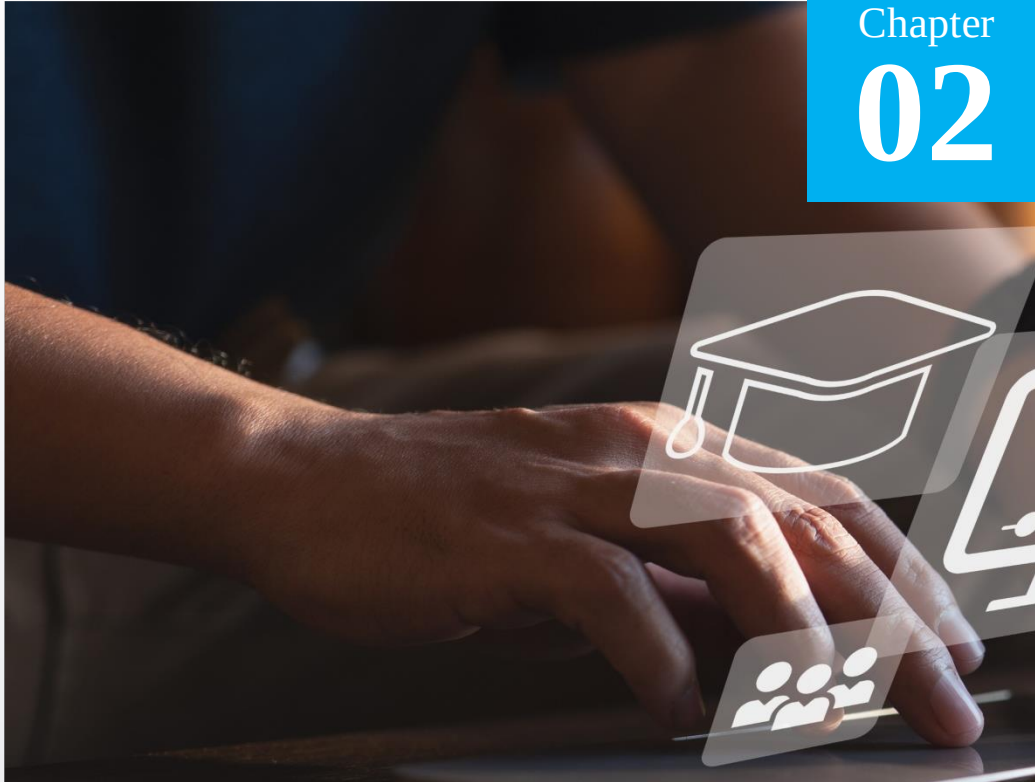
Merge sort:

- Sorting Problem:
 - Sort a sequence of n elements into non decreasing order (For example, 1, 3, 3, 6, 8, 9 **are in non-decreasing order**).
- Divide:
 - Divide the n -element sequence into two subsequences of $n/2$ elements each
- Conquer:
 - Sort the two subsequences recursively using merge sort.
- Combine:
 - Merge the two sorted subsequences to produce the sorted sequence

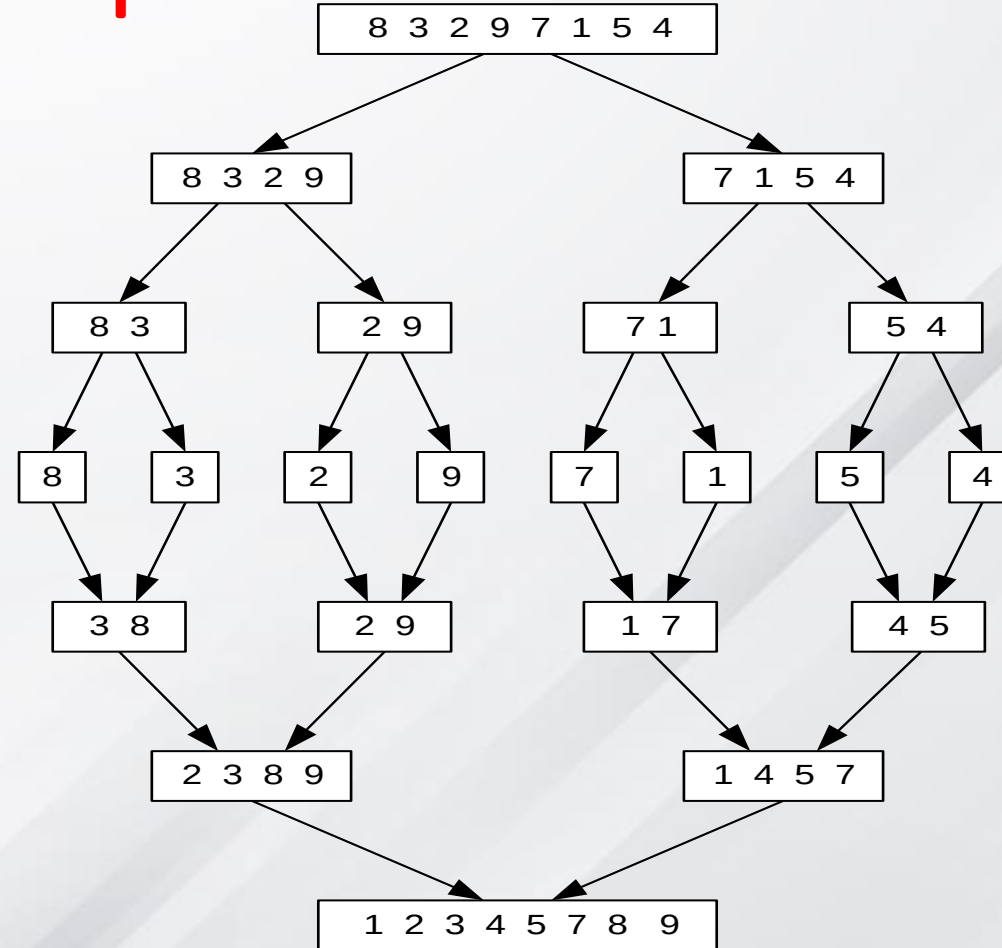
Chapter
02

MERGESORT

Algorithm, Example



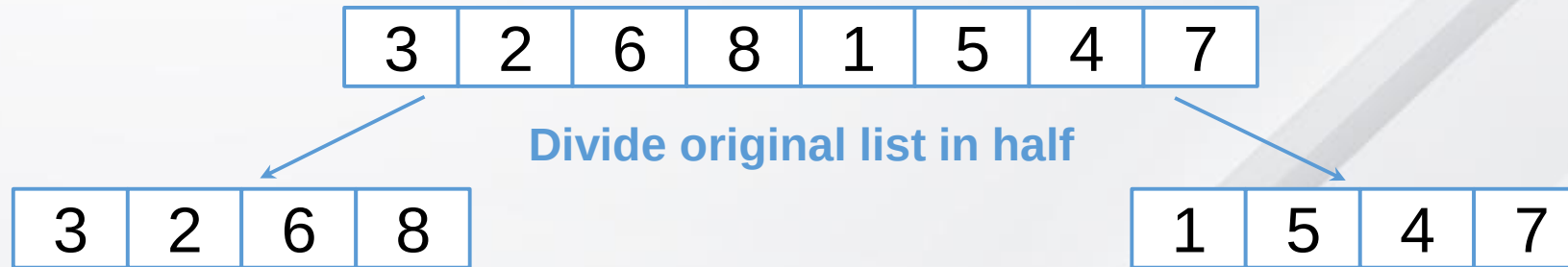
Mergesort Example:



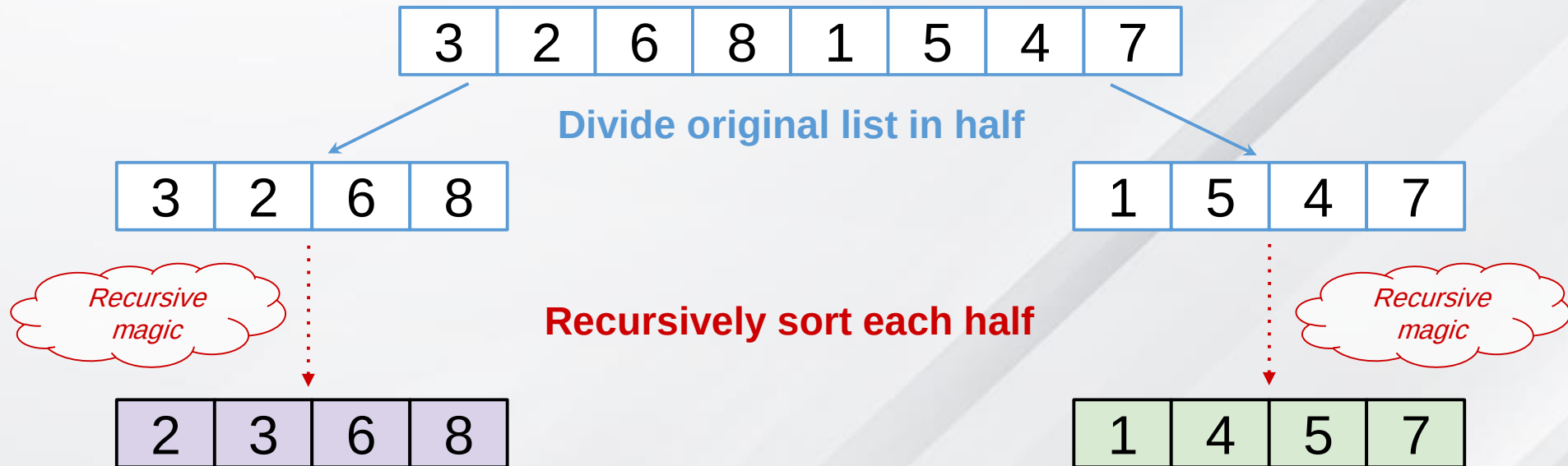
Mergesort:

3	2	6	8	1	5	4	7
---	---	---	---	---	---	---	---

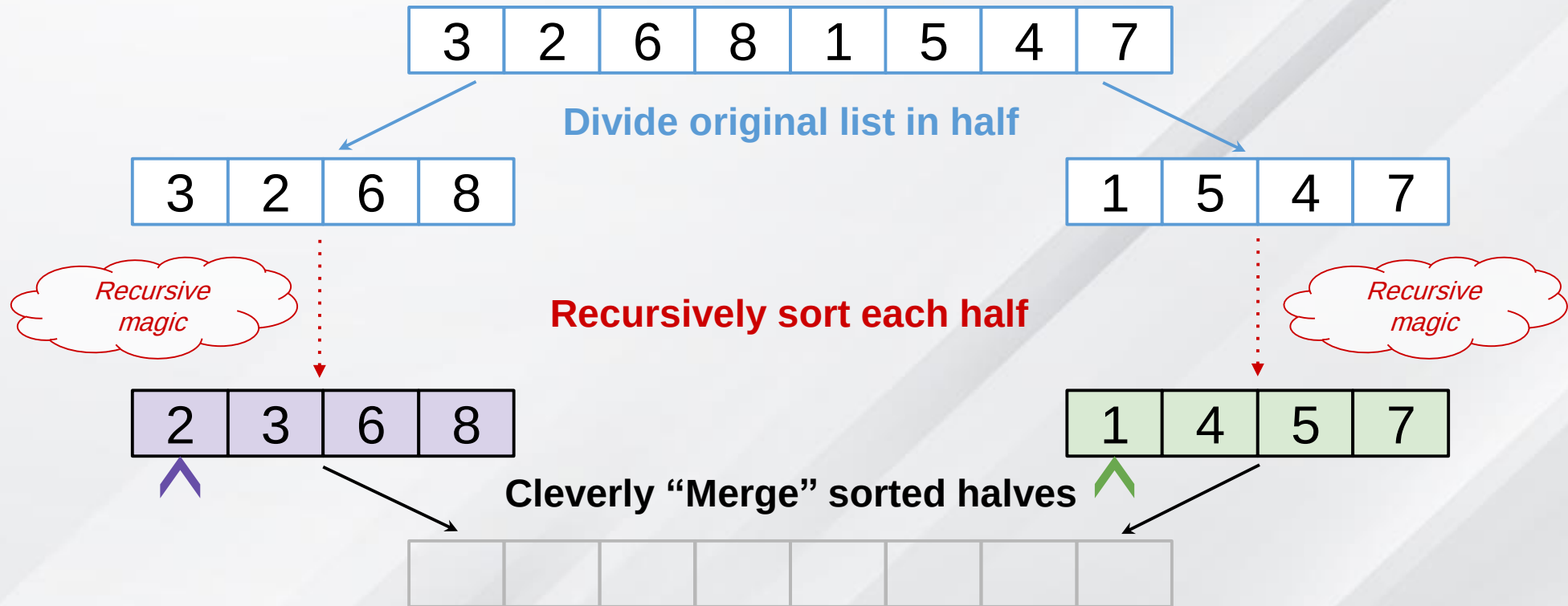
Mergesort:



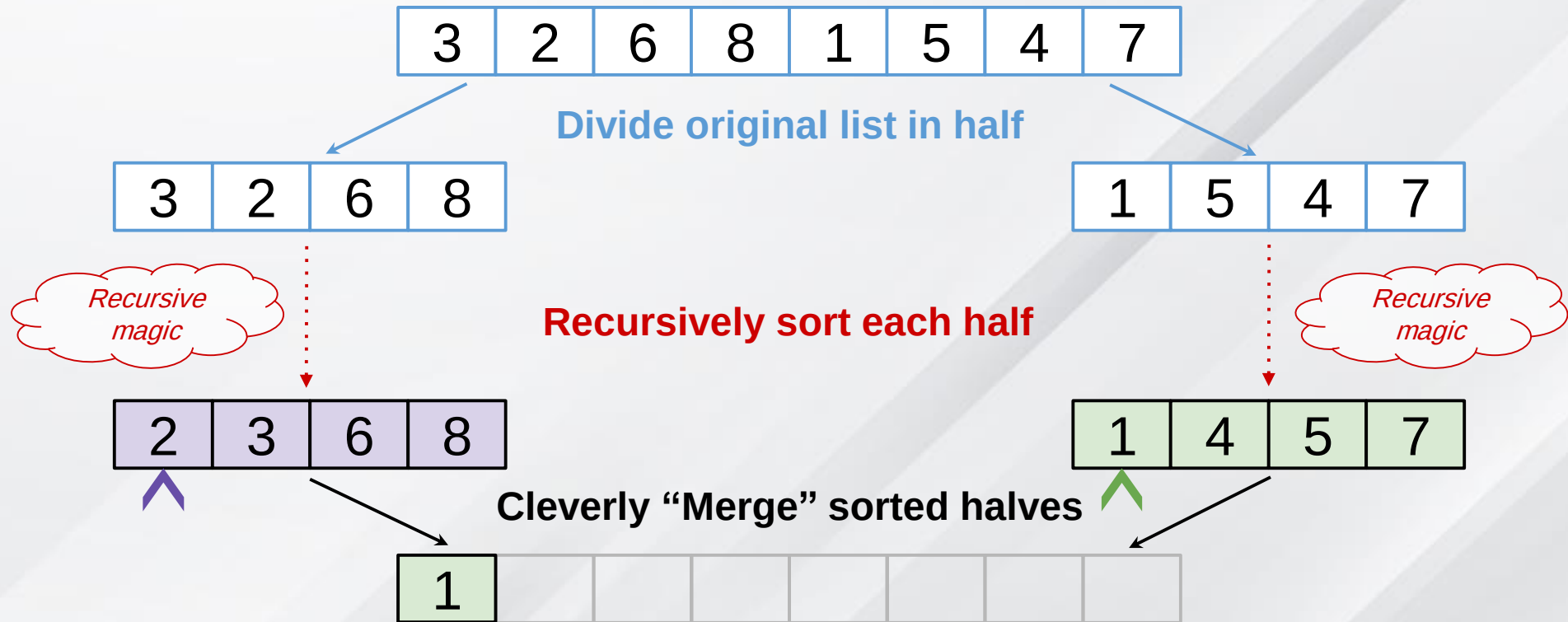
Mergesort:



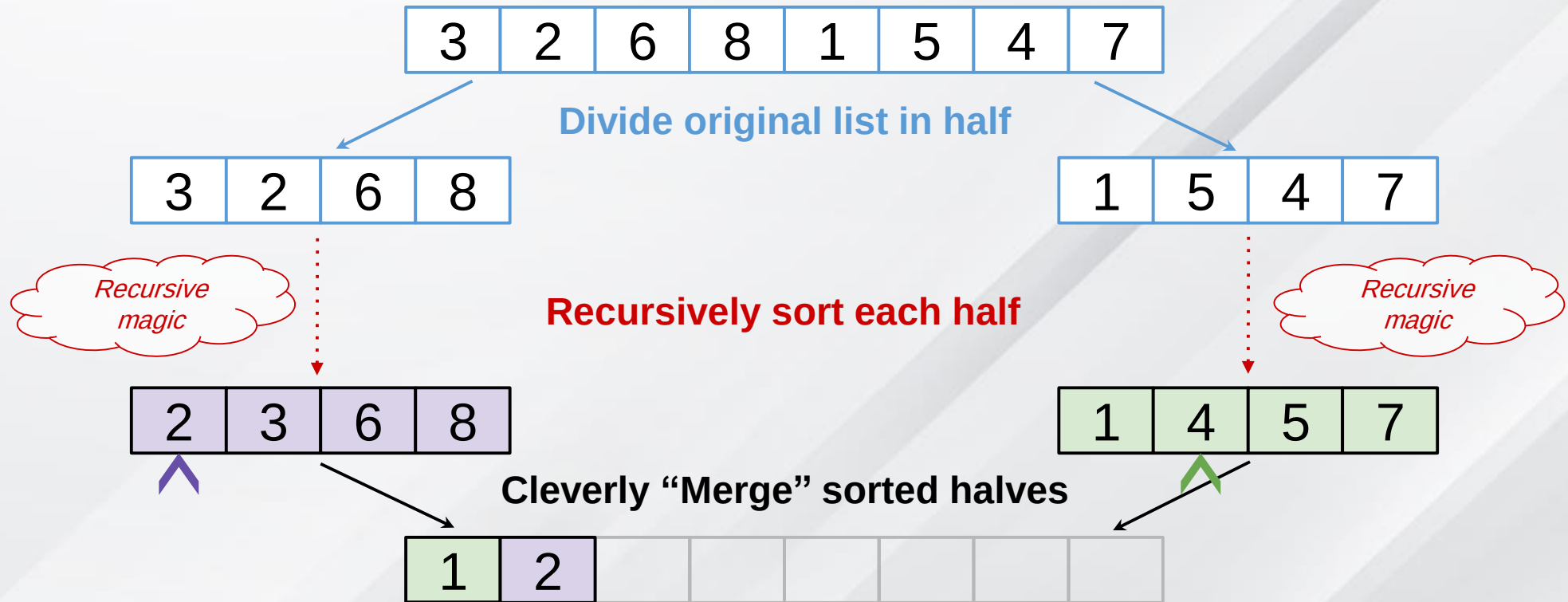
Mergesort:



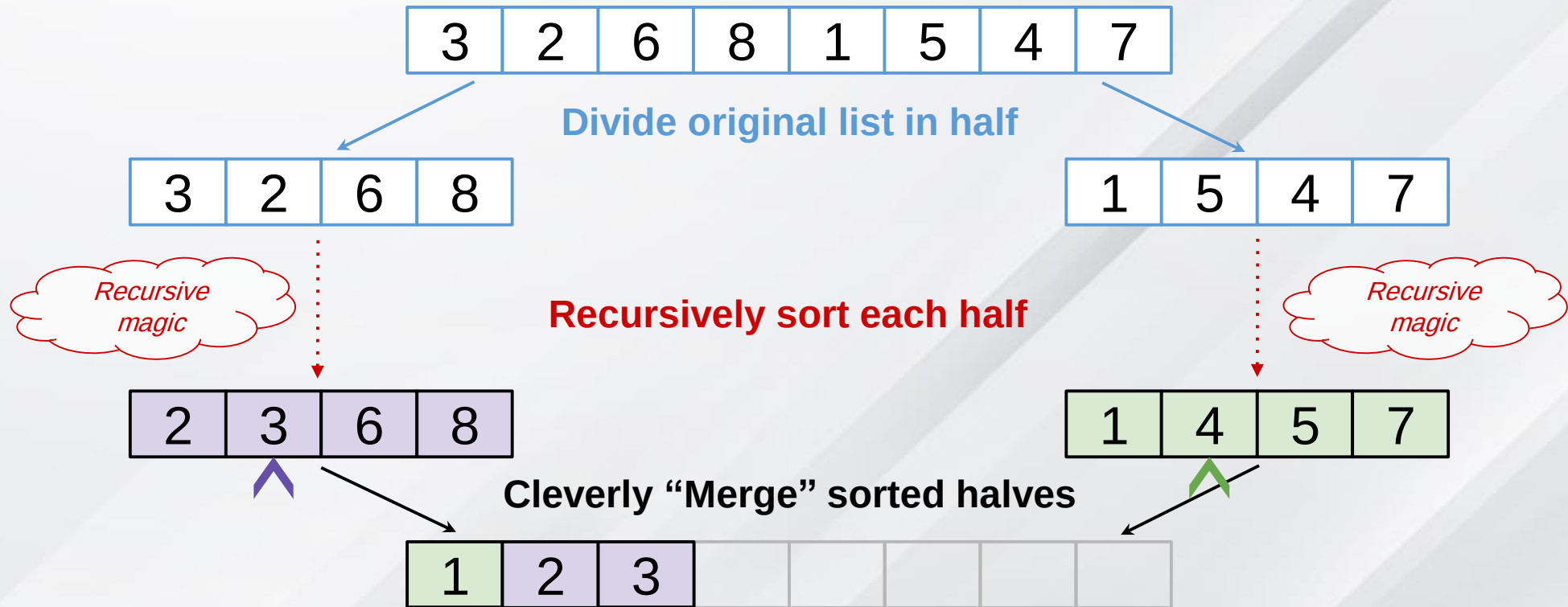
Mergesort:



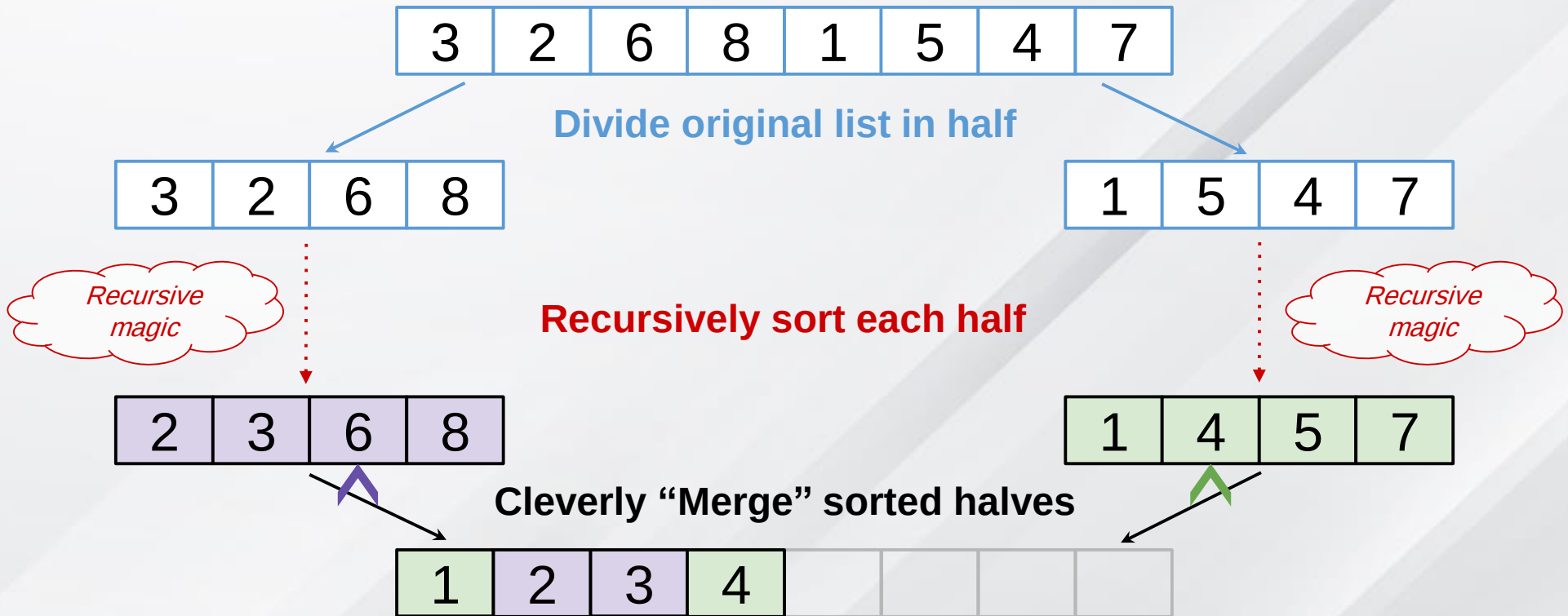
Mergesort:



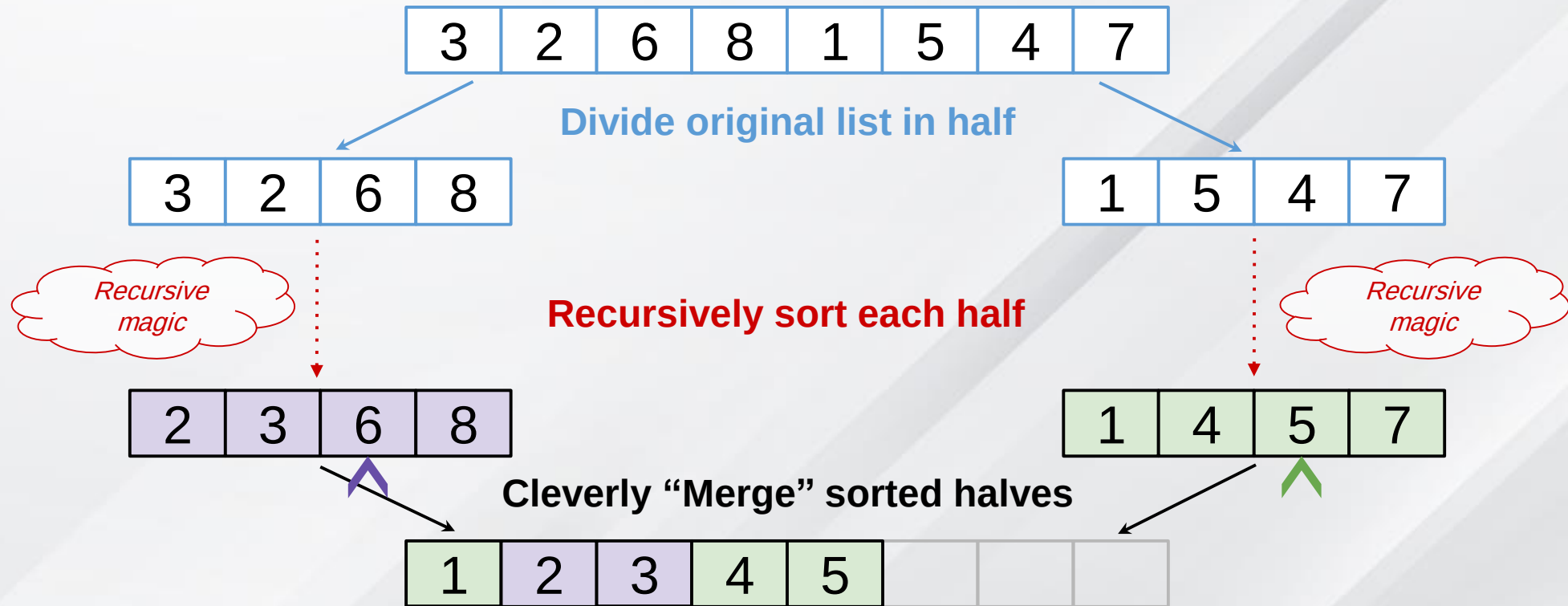
Mergesort:



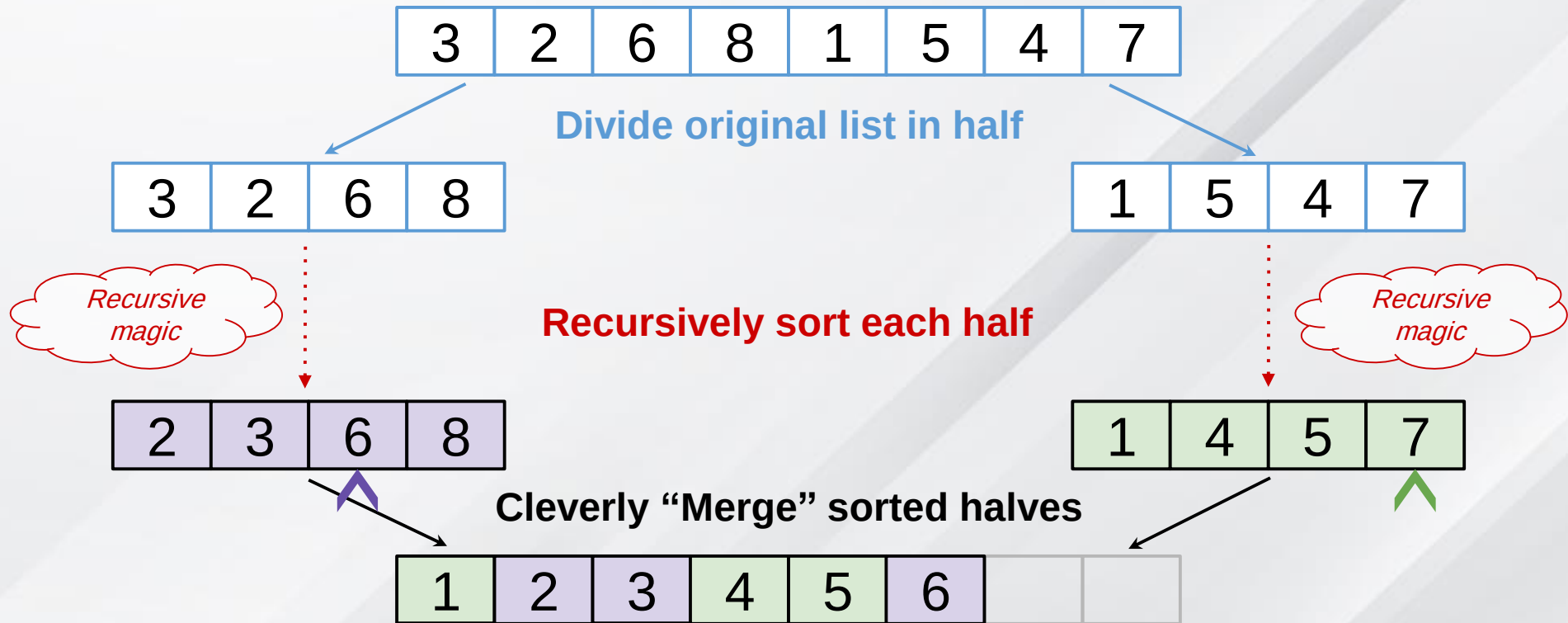
Mergesort



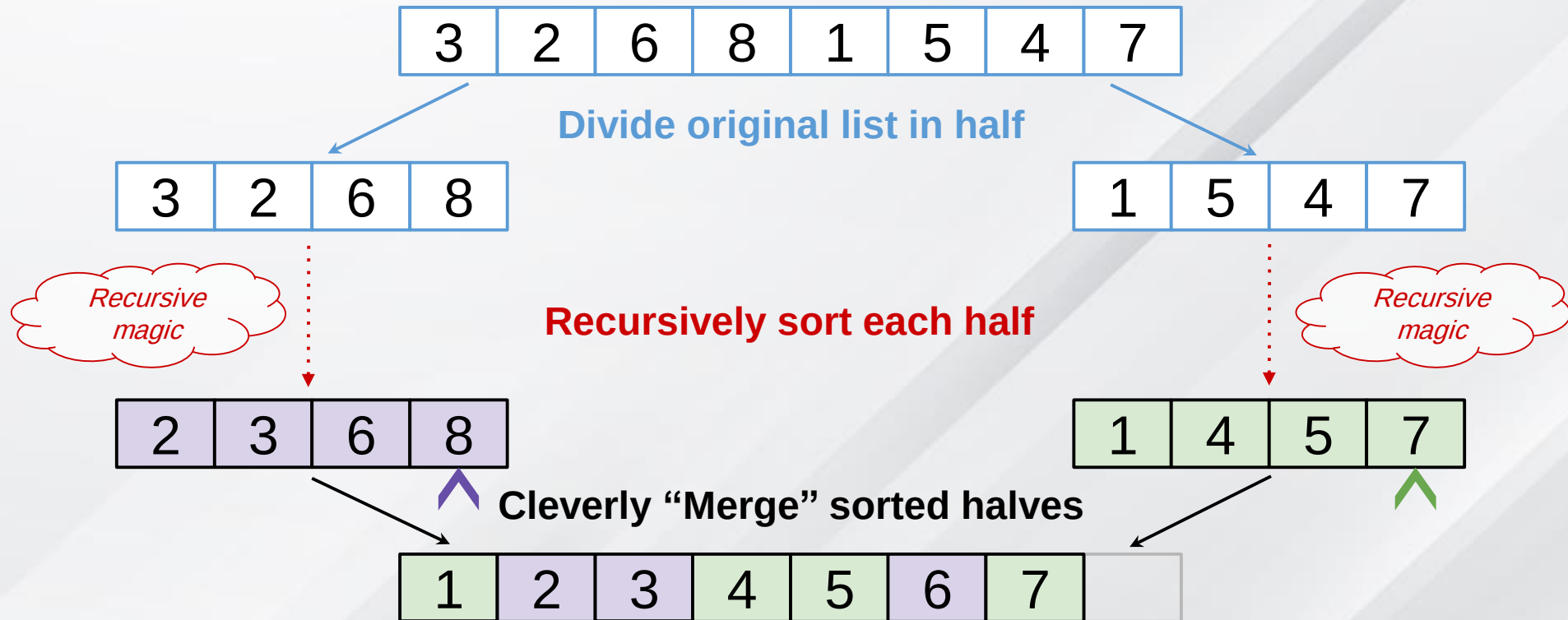
Mergesort



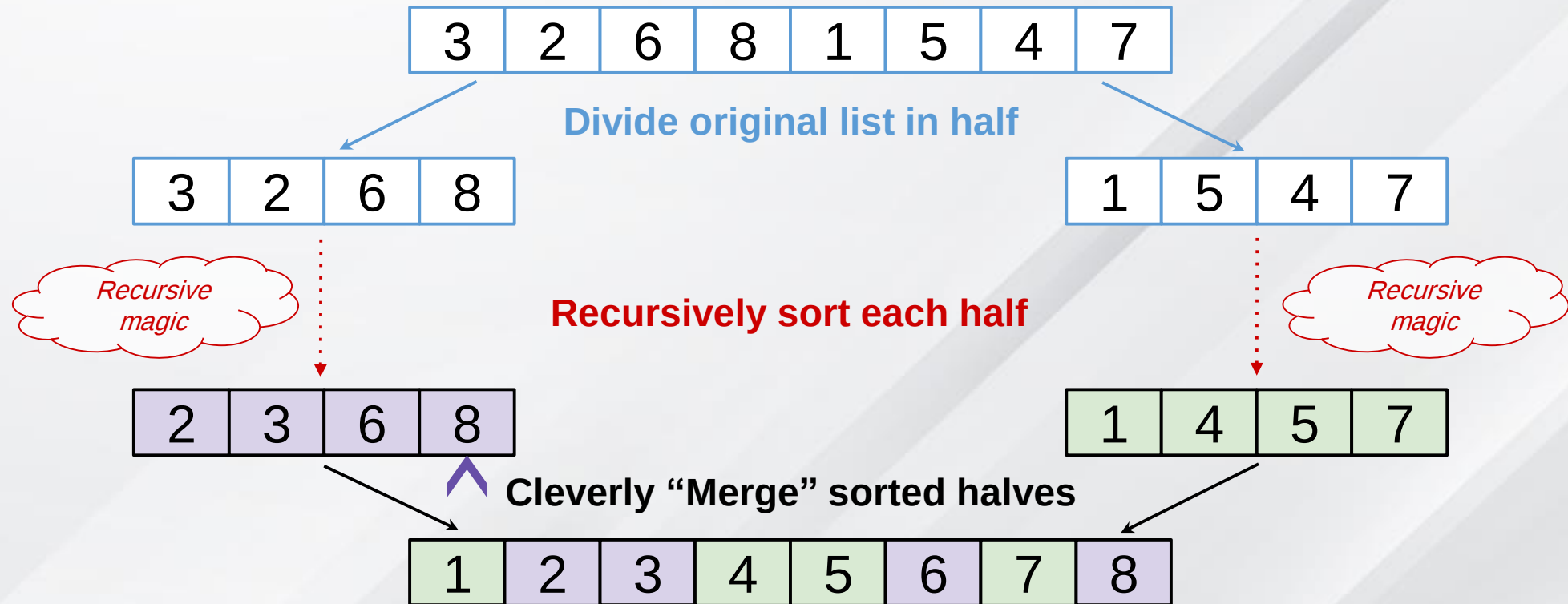
Mergesort



Mergesort



Mergesort



Mergesort: Pseudocode

Intuition: Divide and Conquer. If you sort your left and right halves, it's easier to “Merge” them into a sorted list.

MERGESORT(A):

Mergesort: Pseudocode

Intuition: Divide and Conquer. If you sort your left and right halves, it's easier to “Merge” them into a sorted list.

MERGESORT(A):

$n = \text{len}(A)$

 if $n \leq 1$:

 return A

Mergesort: Pseudocode

Intuition: Divide and Conquer. If you sort your left and right halves, it's easier to “Merge” them into a sorted list.

MERGESORT(A):

$n = \text{len}(A)$

if $n \leq 1$:

 return A

$L = \text{MERGESORT}(A[0:n/2])$

Mergesort: Pseudocode

Intuition: Divide and Conquer. If you sort your left and right halves, it's easier to “Merge” them into a sorted list.

MERGESORT(A):

$n = \text{len}(A)$

 if $n \leq 1$:

 return A

$L = \text{MERGESORT}(A[0:n/2])$

$R = \text{MERGESORT}(A[n/2:n])$

Mergesort: Pseudocode

Intuition: Divide and Conquer. If you sort your left and right halves, it's easier to "Merge" them into a sorted list.

MERGESORT(A):

$n = \text{len}(A)$

 if $n \leq 1$:

 return A

If A has length 1,
It is already sorted!

 L = MERGESORT(A[0:n/2])

 R = MERGESORT(A[n/2:n])

 return MERGE(L,R)

For today, let's assume that n is a power of 2.

Sort the left half

Sort the right half

Merge the two halves

Mergesort: Pseudocode

MERGESORT(A):

$n = \text{len}(A)$

 if $n \leq 1$:

 return A

$L = \text{MERGESORT}(A[0:n/2])$

$R = \text{MERGESORT}(A[n/2:n])$

 return **MERGE**(L,R)

MERGE(L,R):

 result = length n array

$i = 0, j = 0$

 for k in $[0, \dots, n-1]$:

 if $L[i] < R[j]$:

 result[k] = L[i]

$i += 1$

 else:

 result[k] = R[j]

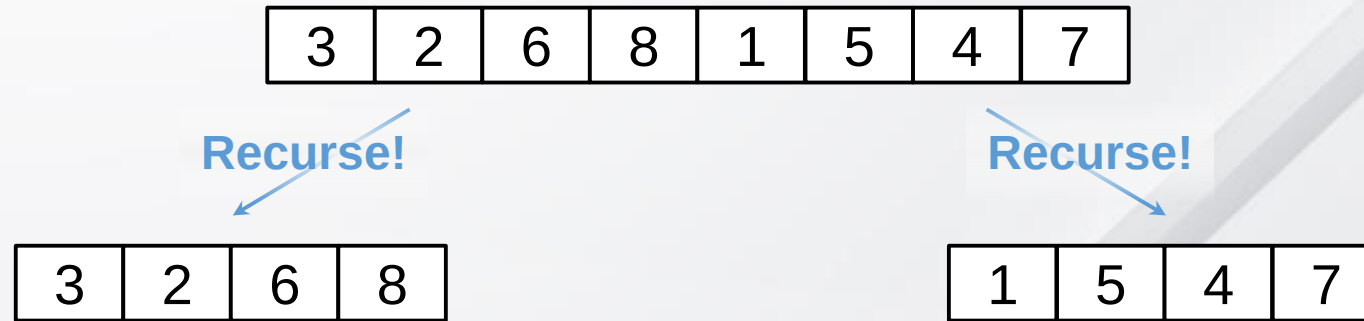
$j += 1$

 return result

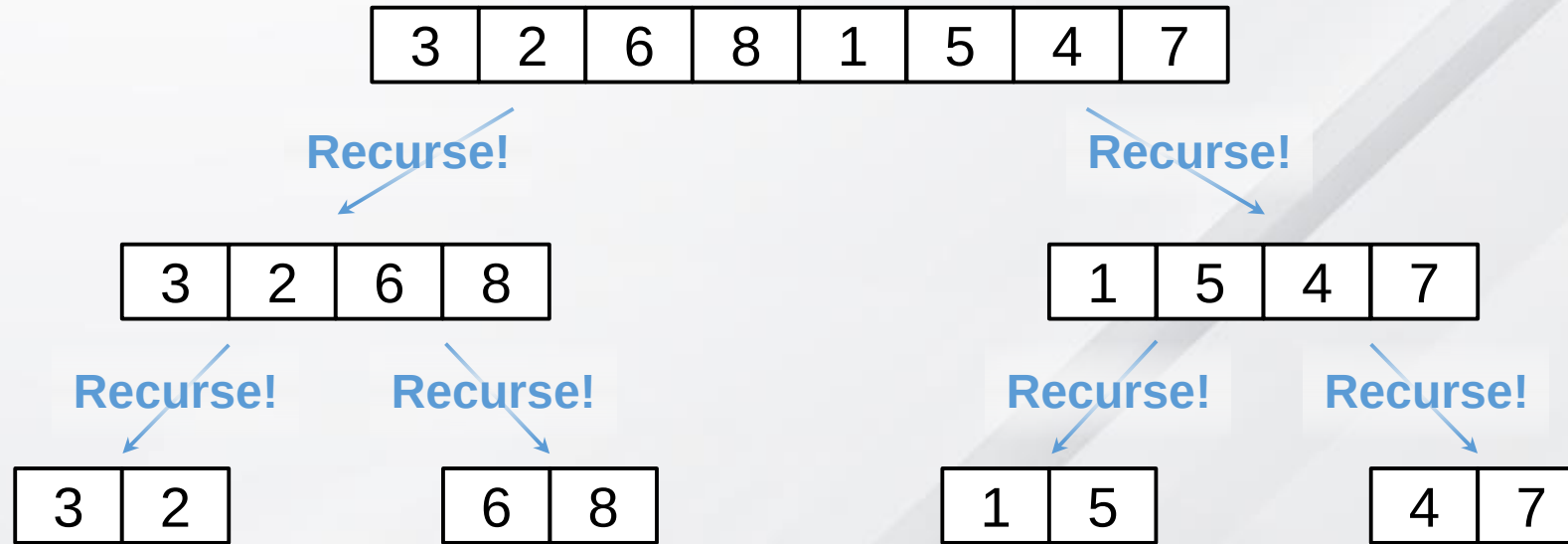
Mergesort: Recursive Calls

3	2	6	8	1	5	4	7
---	---	---	---	---	---	---	---

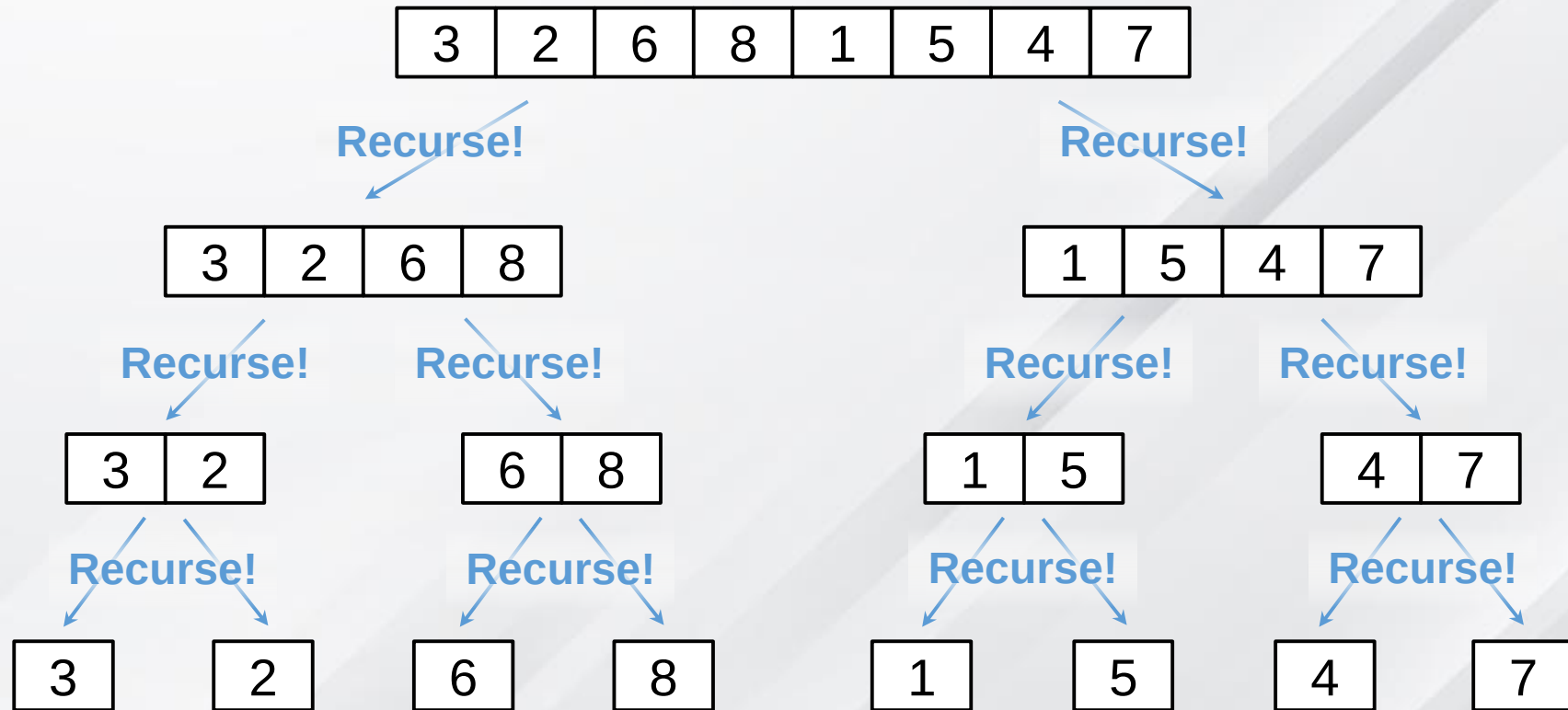
Mergesort: Recursive Calls



Mergesort: Recursive Calls



Mergesort: Recursive Calls

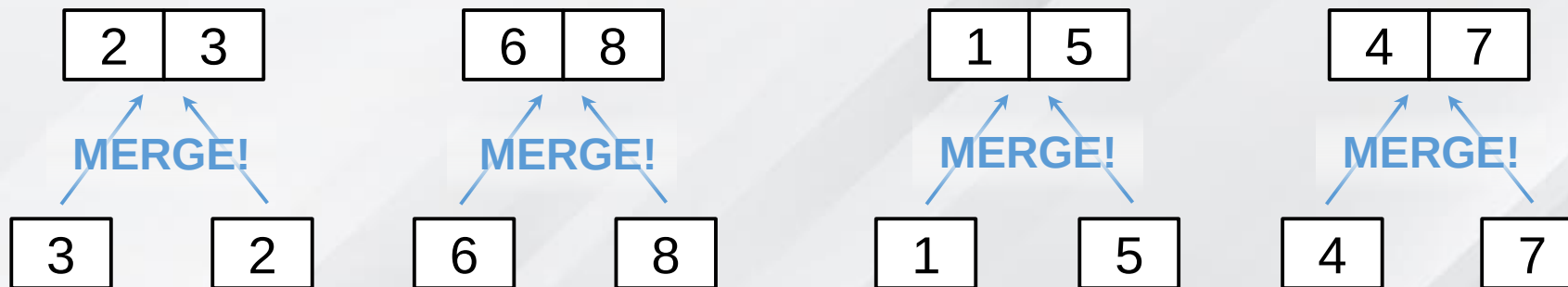


**This is
where we hit
our base
case!**

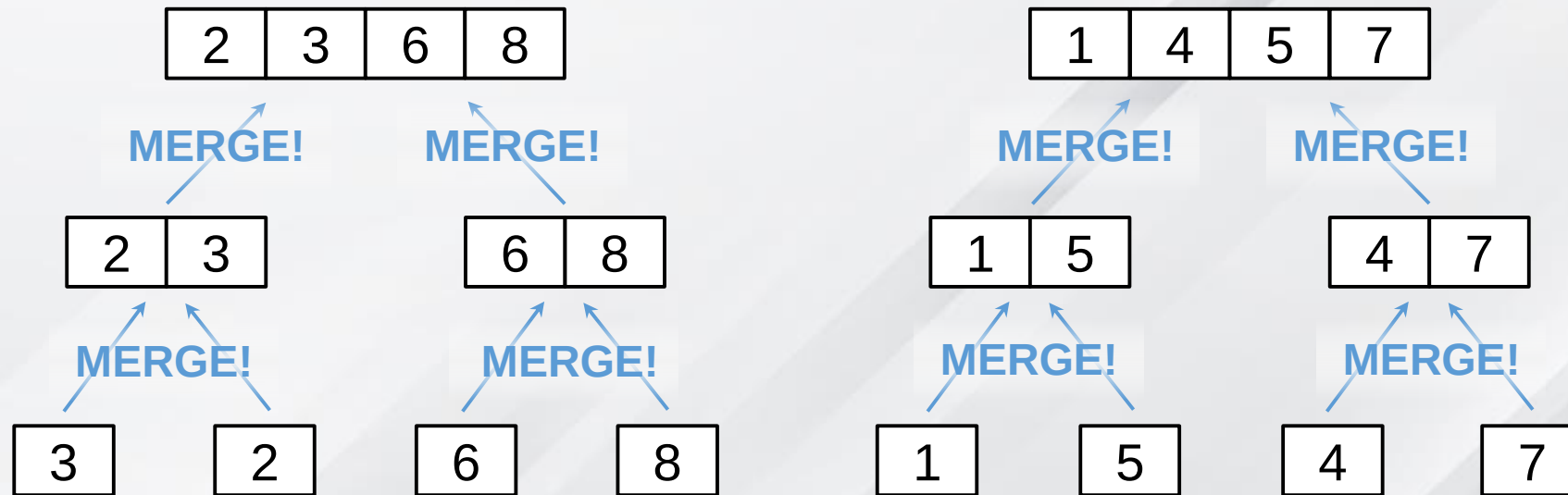
Mergesort: Recursive Calls

3 2 6 8 1 5 4 7

Mergesort: Merge Steps



Mergesort: Merge Steps



Mergesort: Merge Steps



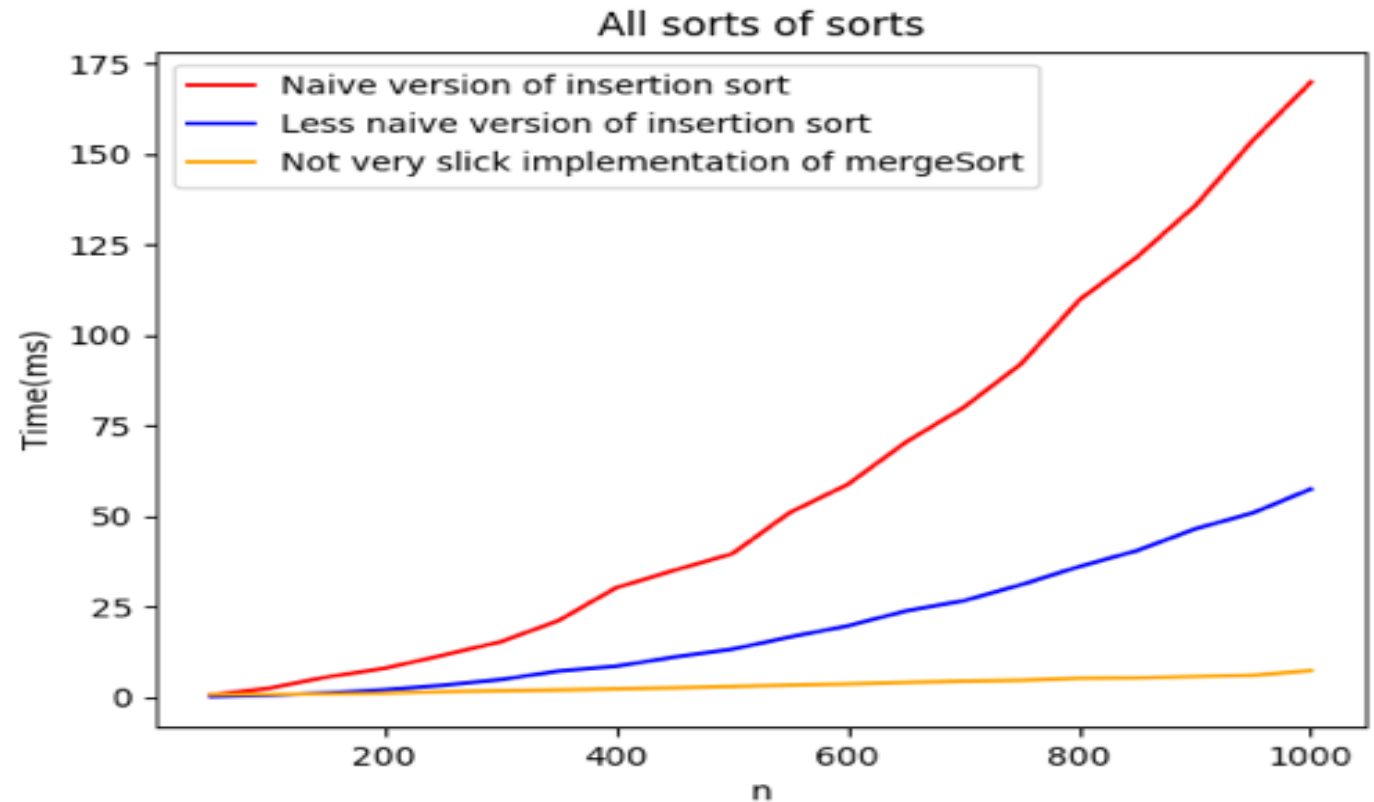
Two questions

1. Does this work?
2. Is it fast?

Empirically:

1. Seems to work.
2. Seems fast.

Python says...



Mergesort: Is It Fast?

MERGESORT(A):

$n = \text{len}(A)$

 if $n \leq 1$:

 return A

$L = \text{MERGESORT}(A[0:n/2])$

$R = \text{MERGESORT}(A[n/2:n])$

 return **MERGE**(L,R)

CLAIM: MergeSort runs in time $O(n \log n)$



Thank You